

Efficiently Learning the Topology and Behavior of a Networked Dynamical System Via Active Queries

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|---|---|
| Dynamical Systems: Basics | |
| A Synchronous Dynamical System (SyDS) has: | |
| • An underlying graph $G(V, E)$. | |
| • Nodes: Agents in the system. | |
| • Edges: Permissible local interactions. | |
| • State values for nodes from a finite | |
| domain (here: {0, 1}). A local transition function for each node. | |
| Update mechanism: synchronous. | |
| Local Transition Function: | |
| | |
| $\begin{array}{cccc} a & b & c & Local function f_v: \\ \hline & & & \\ \end{array}$ • Inputs: States of nodes | |
| 1100000000000000000000000000000000000 | |
| • - state 0 hood $N^+(v)$ of node v . | |
| • - state 1 \bigvee_{v} • <u>Output</u> : Next state of v . | |
| Symmetric Function: | |
| • Each node v has a table T_v with $deg_v + 1$ rows. | |
| • When the number of 1's in the input is i , the | |
| next state of v is given by $T_v[i]$. Threshold Function: | |
| • Each node v has an integer threshold τ_v . | |
| • The next state of v is 1 iff the no. of state-1 | |
| nodes in $N^+(v)$ is at least τ_v . | |
| Example of a Threshold-SyDS: | |
| $oldsymbol{O}$ - state 0 $ullet$ - state 1 $	au_a=2$ $	au_b=1$ $	au_c=2$ $	au_d=2$ | |
| a b c a b c a b c | |
| | |
| $\sum_{i=1}^{n} \sum_{i=1}^{n} \sum_{i$ | |
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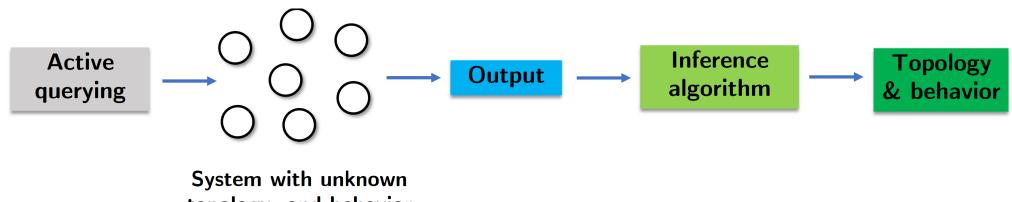
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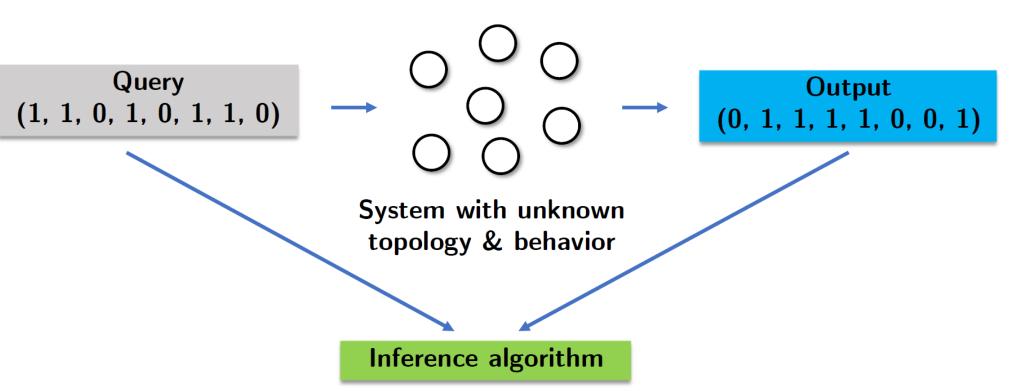
PROBLEM OVERVIEW

Our focus: Inferring both the topology and the local functions of a SyDS.



Query Model

- Each query specifies the *input* to each of the local functions. The response provides the value of each function for the inputs specified by the query.
- **Batch mode**: Submit all queries in a *single batch*.
- Adaptive mode: Submit Queries in *multiple* stages; new queries may be based on the responses to the previous queries.



PREVIOUS WORK (BRIEF)

- Active querying to predict users' choices from a known set of options [Kleinberg et al., 2017].
- Inferring influence functions for networked systems [Romero et al., 2011, Narasimhan et al., 2015].
- Inferring the network structure given the contagion propagation model [Abrahao et al., 2013, Gomez-Rodriguez et al., 2010].
- Learning the local functions of discrete dynamical systems from *observed* data [Adiga et al., 2017].
- Given the network topology, active querying to learn the local functions [Adiga et al., 2018; 2019; 2020].

CONTACT

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Theorem 2 (Adaptive: $O(n + m \log n)$ **queries)** In polynomial time, an *adaptive* query set Q with $O(n + m \log n)$ queries can be constructed, such that *Q* can be used to identify the local functions and the topology of a symmetric-SyDS S.

Remarks:

To infer threshold-SyDSs: A deterministic algorithm under the adaptive mode and a randomized algorithm under the batch mode.

Theorem 3 (Adaptive: $O(n + m \log n)$ **queries)** In polynomial time, an *adaptive* query set Q with $O(n + m \log n)$ queries can be constructed, such that *Q* can be used to identify the local functions and the topology of a threshold-SyDS S.

Theorem 4 (Batch: Randomized) In polynomial time, a *batch* query set Q with $O(n\Delta \log n)$ queries can be constructed, such that *Q* can be used to identify the local functions and the topology of a threshold-SyDS S w.h.p.

Remark: The randomized algorithm assumes that an upper bound on the maximum degree Δ is known.

OUR CONTRIBUTIONS

To infer symmetric-SyDSs: Efficient inference algorithms under batch and adaptive modes.

Theorem 1 (Batch: $O(n^2)$ **queries)** In polynomial time, a *batch* query set Q with $O(n^2)$ queries can be constructed, such that Q can be used to identify the local functions and the topology of S.

• No. of nodes = n and number of edges = m.

• The algorithms don't assume that *m* is known; the parameter m in the bound results from the analysis. • For *sparse* networks (e.g., m = O(n)), the adaptive mode algorithm uses asymptotically fewer queries.

A lower bound for batch mode:

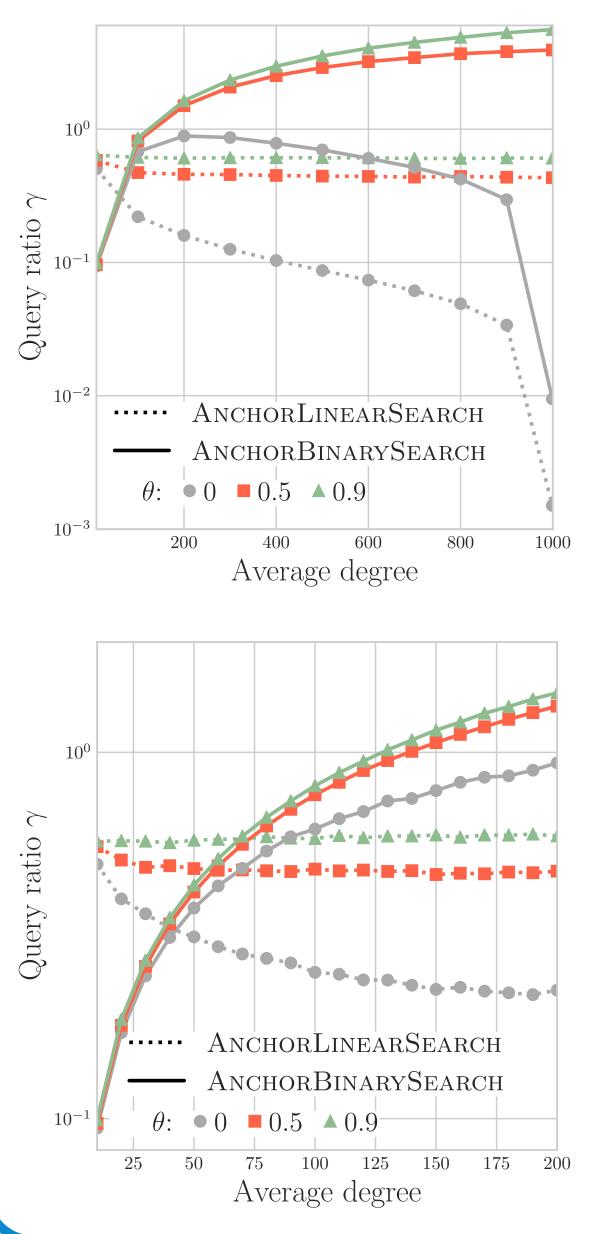
Theorem 5 (Batch: Lower bound) Under the batch mode, any query set that can correctly learn the network topology of any threshold-SyDS must contain $\Omega(n \log n)$ queries.

EXPERIMENTAL RESULTS

Experimental Setup:

- law networks.

Experimental Results (Examples):





OUR CONTRIBUTIONS (CONT.)

• **Datasets**: (*i*) real-world networks from various

domains, with the numbers of nodes ranging from 7,000+ to 28,000+; (ii) synthetic gnp and power-

• Metric: Query ratio $\gamma = a/b$, where a and b are the numbers of queries used by an adaptive algorithm and a batch algorithm respectively.

> Gnp network with 10,000 nodes and average degree varied from 10 to 999.

Power-law network with 10,000 nodes and average degree varied from 10 to 200.